Applied Compositional Thinking for Engineers



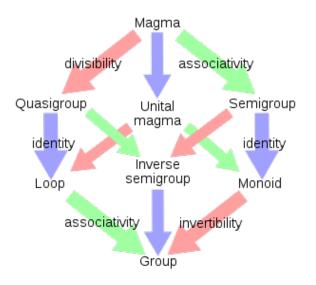
Spring 2021

Semicategories, Categories

Building in layers

• Last week:

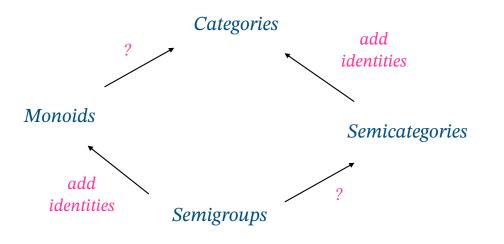
- A magma is a set with a binary operation of *composition*.
- A **semigroup** is a magma whose composition is **associative**.
- A monoid is a semigroup with a neutral element.
- A group has is a monoid with an "inverse" operation.





Building in layers

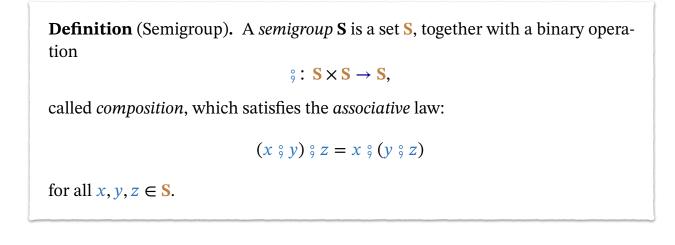
• Today: we will *generalize*, rather than *refine*.





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1 8	11.07	Romanshorn	34	
S8	11.07	Pfäffikon SZ	32	
S9	11.07	Schaffhausen	41/42	
//र 37	11.08	Basel SBB	13	
EC	11.09	Como S. Giovanni	9	
//र 37	11.09	St. Gallen	11	
S5	11.09	Zug	41/42	
//₹ 36	11.10	Basel SBB	31	
S15	11.10	Rapperswil	43/44	
S19	11.11	Dietikon	32	+1'





• Examples:

 $\langle \mathbb{R}, \cdot
angle$

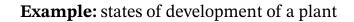
 $\langle \mathbb{R}, +
angle$

⟨ℕ, min⟩

 $\langle \mathbb{N}, \max \rangle$

 $\langle \mathbb{R}, \max \rangle$

. 2



 $X = \{$ sprout, young, mature, old, dead $\}$

 $T: X \to X$

T(sprout) = youngT(young) = matureT(mature) = oldT(old) = deadT(dead) = dead

$$\mathbf{S} = \{T^n \mid n \in \mathbb{N}\} \qquad \qquad T_j^* \vdash = \top, \quad T_j^* \vdash = \top, \quad \dots$$

 $T^n = T^4 \quad \forall \ n \geq 4$

 $\mathbf{S} = \{T, T^2, T^3, T^4\}$

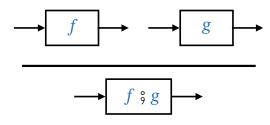


Example: discrete time linear time-invariant systems of the form:

$$x_{k+1} = \mathbf{A}x_k + \mathbf{B}u_k$$
$$y_k = \mathbf{C}x_k$$

with the constraint that input and output have the same dimension.

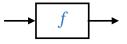
- The composition is the series composition.
- The composition of linear system is linear
 → we have a semigroup.



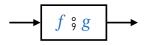


Example: discrete time linear time-invariant systems of the form:

$$x_{k+1} = \mathbf{A}x_k + \mathbf{B}u_k$$
$$y_k = \mathbf{C}x_k$$









Definition 11.6 (Semicategory). A semicategory C is:

Constituents

- 1. Objects: a collection[‡] Ob_C, whose elements are called *objects*.
- 2. Morphisms: for every pair of objects $X, Y \in Ob_{\mathbb{C}}$, there is a set $\operatorname{Hom}_{\mathbb{C}}(X; Y)$, elements of which are called *morphisms* from X to Y. The set is called the "hom-set from X to Y".
- 3. Composition operations: given any morphism $f \in \text{Hom}_{\mathbb{C}}(X; Y)$ and any morphism $g \in \text{Hom}_{\mathbb{C}}(Y; Z)$, there exists a morphism $f \ ; g \in \text{Hom}_{\mathbb{C}}(X; Z)$ which is the *composition of f and g*.

Conditions

1. Associativity: for any morphisms $f \in \text{Hom}_{\mathbb{C}}(X; Y), g \in \text{Hom}_{\mathbb{C}}(Y; Z)$, and $h \in \text{Hom}_{\mathbb{C}}(Z; W)$,

$$(f \ ; g) \ ; h = f \ ; (g \ ; h).$$
 (11.1)



Example: discrete time linear time-invariant systems of the form:

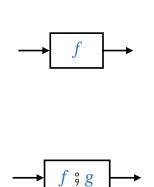
 $x_{k+1} = \mathbf{A}x_k + \mathbf{B}u_k$ $y_k = \mathbf{C}x_k$

• Objects:

• Morphisms:

• Composition:

• Associativity:





Example: states of a plant

 $X = \{$ sprout, young, mature, old, dead $\}$

 $T: X \to X$

T(sprout) = youngT(young) = matureT(mature) = oldT(old) = deadT(dead) = dead

 $\mathbf{S} = \{T, T^2, T^3, T^4\}$

 $R : X \rightarrow Y$ R(sprout) = alive R(young) = alive R(mature) = alive R(old) = alive R(dead) = dead

 $Y = \{\text{alive, dead}\}$ $T' : Y \to Y$

T'(alive) = deadT'(dead) = dead

$$\mathbf{S}' = \{T'\}$$

Example: states of a plant

$X = \{$ sprout, young, mature, old, dead $\}$	$Y = \{alive, dead\}$	
$T: X \to X$	$T': Y \to Y$	
T(sprout) = young		
T(young) = mature	T'(alive) = dead T'(dead) = dead	
T(mature) = old		
T(old) = dead		
T(dead) = dead		
$\mathbf{S} = \{T, T^2, T^3, T^4\}$	$\mathbf{S}' = \{T'\}$	

 $R:X\to Y$







A monoid is a semigroup with a *neutral element*.

Definition (Monoid). A monoid **M** is: <u>Constituents</u> 1. a set **M**; 2. a binary operation $\S: \mathbf{M} \times \mathbf{M} \to \mathbf{M}$; *new* 3. a specified element id $\in \mathbf{M}$, called *neutral element*. <u>Conditions</u> 1. Associative law: $(x \S y) \S z = x \S (y \S z)$; *new* 2. Neutrality Laws: id $\S x = x = x \S$ id.

Examples

 $\langle \mathbb{R}, \cdot
angle$

 $\langle \mathbb{R}, + \rangle$

⟨ℕ, min⟩

⟨ℕ, max⟩

 $\langle \mathbb{R}, \max \rangle$



- Sometime we can "formally add" a neutral element to a semigroup.
- Example:

 $\langle \mathbb{N}, \min \rangle$



- Sometime we can "formally add" a neutral element to a semigroup.
- Example:
 - $\langle \mathbb{R}, \max \rangle$



• Sometime we can "formally add" a neutral element to a semigroup.

• Example:

 $\mathbf{A} = \{\bullet, \bullet\}$

 $S = \{ set of non-empty strings of elements of A \}$



Example: states of development of a plant

 $X = \{$ sprout, young, mature, old, dead $\}$

 $T: X \to X$

 $\mathbf{S}=\{T,T^2,T^3,T^4\}$

 $\mathbf{M} = \{ \mathrm{id}, T, T^2, T^3, T^4 \}$



From semicategories to categories

Example: states of development of a plant

$X = \{$ sprout, young, mature, old, dead $\}$	$Y = \{alive, dead\}$
$T: X \to X$	$T':Y\to Y$
$\mathbf{S} = \{T, T^2, T^3, T^4\}$	$\mathbf{S'} = \{T'\}$

 $R:X\to Y$



From semicategories to categories

Example: states of development of a plant

$X = \{$ sprout, young, mature, old, dead $\}$	$Y = \{\text{alive, dead}\}\$
$T: X \to X$	$T':Y\to Y$
$\mathbf{M} = \{ \mathrm{id}, T, T^2, T^3, T^4 \}$	$\mathbf{M'} = \{\mathrm{id}, T'\}$

 $R:X\to Y$



Definition 1.10 (Category). A category C is:

Constituents

- Objects: a collection Ob_c, whose elements are called *objects*.
- Morphisms: for every pair of objects X, Y ∈ Ob_C, there is a set Hom_C (X; Y), elements of which are called *morphisms* from X to Y. The set is called the "hom-set from X to Y".
- *new* 3. Identity morphisms: for each object X, there is an element $Id_X \in Hom_{\mathbb{C}}(X;X)$ which is called *the identity morphism of X*.
 - 4. Composition operations: given any morphism $f \in \operatorname{Hom}_{\mathbb{C}}(X;Y)$ and any morphism $g \in \operatorname{Hom}_{\mathbb{C}}(Y;Z)$, there exists a morphism $f \ g \in \operatorname{Hom}_{\mathbb{C}}(X;Z)$ which is the *composition of f and g*.

Conditions

new 1. Unitality: for any morphism $f \in \text{Hom}_{\mathbb{C}}(X; Y)$,

$$\mathrm{Id}_{\mathbf{X}} \ ; \ f = f = f \ ; \ \mathrm{Id}_{\mathbf{Y}} . \tag{1.5}$$

2. Associativity: for any morphisms $f \in \text{Hom}_{\mathbb{C}}(X; Y), g \in \text{Hom}_{\mathbb{C}}(Y; Z)$, and $h \in \text{Hom}_{\mathbb{C}}(Z; W)$,

$$(f \ ; g) \ ; h = f \ ; (g \ ; h).$$
 (1.6)



Example: states of development of a plant

$X = \{$ sprout, young, mature, old, dead $\}$	$Y = \{alive, dead\}$
$T: X \to X$	$T':Y\to Y$
$\mathbf{M} = \{ \mathrm{id}, T, T^2, T^3, T^4 \}$	$\mathbf{M'} = \{\mathrm{id}, T'\}$

 $R:X\to Y$



Example: sets and functions

Example: *finite* sets and functions

Example: *a specified collection of* sets and functions



Example: matrices over the real numbers

Example: finite-dimensional vector spaces over the real numbers

Example: all vector spaces over the real numbers



Example: monoids and their morphisms

A monoid as a category

Given a monoid

 $\langle \mathbf{M}, \mathbf{\hat{s}}, \mathbf{id} \rangle$

we can think of it as a special case of a category.



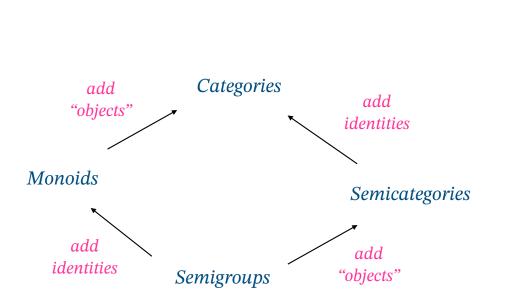
Level shifts

A specific monoid, viewed as a category

 $\langle \mathbf{M}, \mathbf{\hat{g}}, \mathbf{id} \rangle$

The category of all monoids





Summary so far

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